

Geodesic Nets

Construction and Existence

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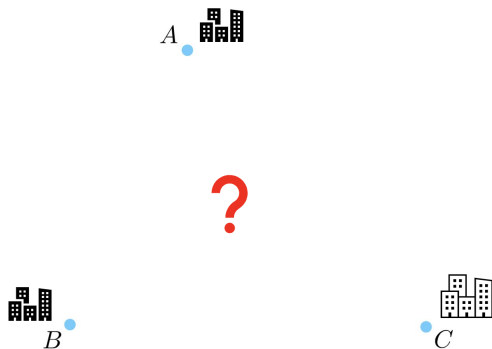
²Texas Christian University, Fort Worth, TX



JOHN V. ROACH
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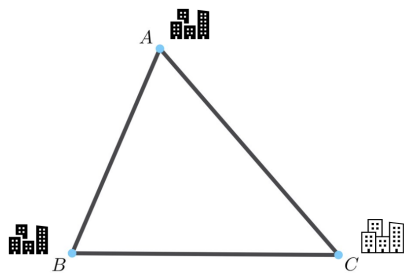
Motivation



Problem: Build a road system such that

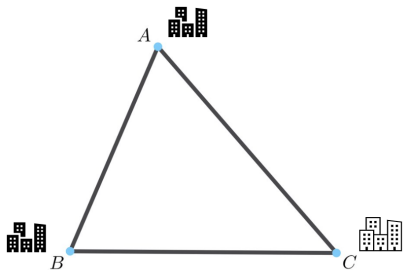
- 1 connecting all three cities
- 2 minimizing the total length of roads

Motivation

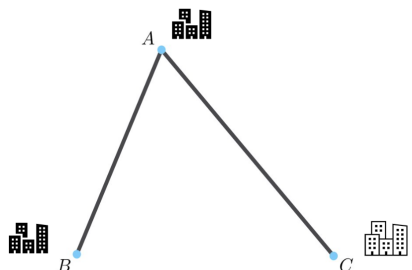


(a) Connect all cities

Motivation

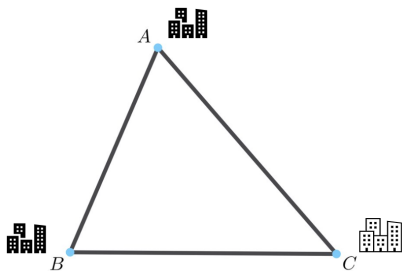


(a) Connect all cities

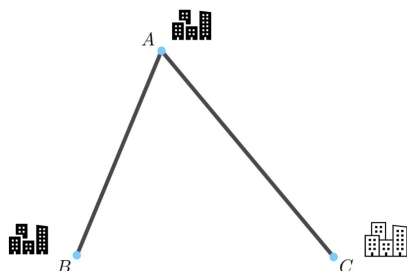


(b) Connect all cities + Smaller total length

Motivation



(a) Connect all cities



(b) Connect all cities + Smaller total length

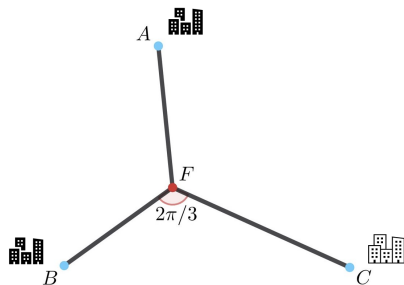
But does this network **minimize** the total length?

Fermat Point

The **Fermat point** F of triangle ABC is the point that **minimizes** the **sum of three distances** from three vertices to the point.

Fermat Point

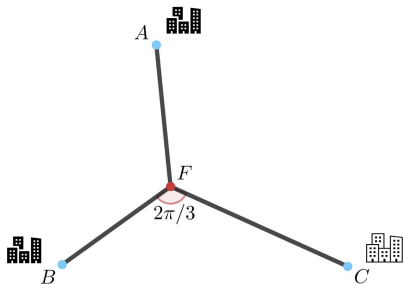
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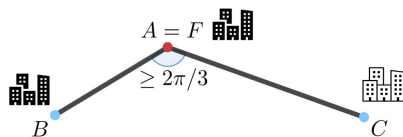
(a) Case 1: all angles $< 2\pi/3$

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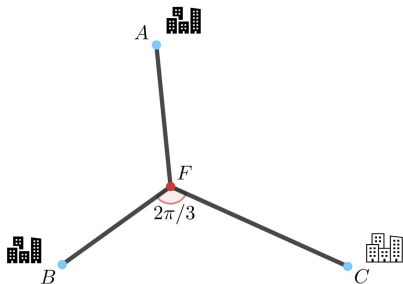
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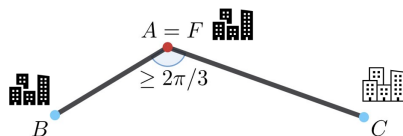
(b) Case 2: $\angle BAC \geq 2\pi/3$

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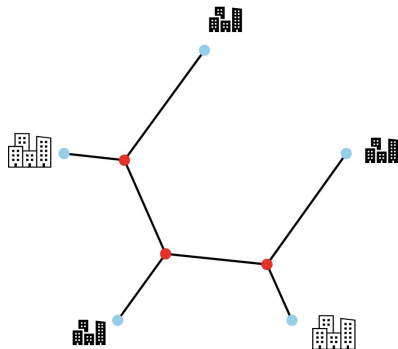
(a) Case 1: all angles $< 2\pi/3$



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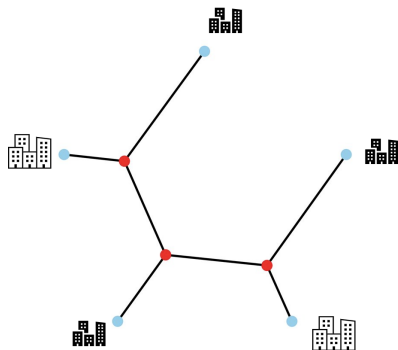
Important property: in Case 1, $\angle AFB = \angle BFC = \angle CFA = 2\pi/3$.

General problems

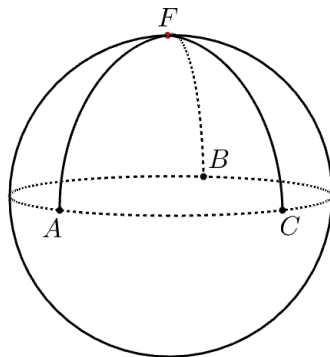


(a) **Steiner Tree** [2] of n points

General problems



(a) **Steiner Tree** [2] of n points



(b) Networks on **different surfaces**

Applications

- **Telecommunication:** design optimal network transporting data between centers
- **Molecular Biology:** find the biological network to investigate the interaction among proteins and genes [1]
- **Urban Planning:** determine the most efficient routes for public transportation systems

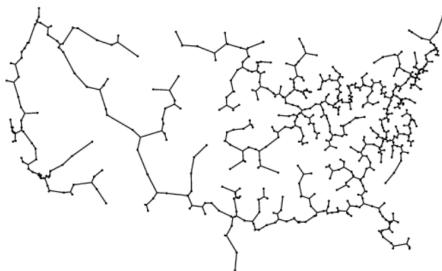
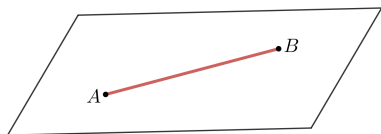


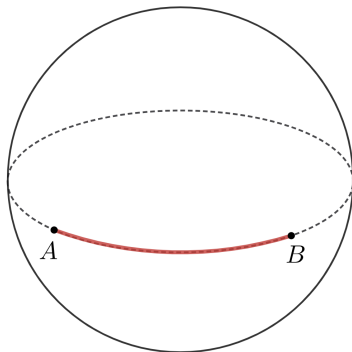
Figure: Steiner Tree of cities in the US [2].

Geodesics

A **geodesic** is a “straight” curve on a surface.



(a) Straight line on a plane



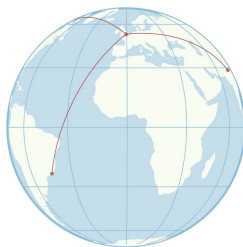
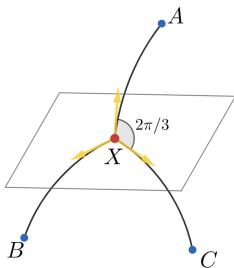
(b) Geodesic on a sphere

Important property: Geodesic **locally** minimizes length.

Geodesic Nets

A **geodesic net** [3] is a **network** connecting points with **geodesics** on surfaces that includes:

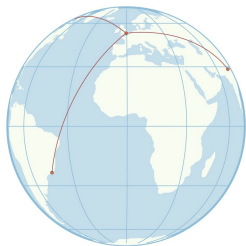
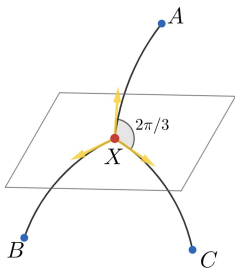
- “unbalanced” points.
- “balanced” points - stretched equally.



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A **geodesic net** [3] is a **network** connecting points with **geodesics** on surfaces that includes:

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- “balanced” points - stretched equally.



Question: How to **construct** geodesic nets given **any number of points**?

Big picture

Main idea: Given n unbalanced points, add new balanced points heuristically.

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- 1 If v_u has degree 1, it remains unchanged.
- 2 If v_u has a degree greater than 1, the algorithm checks whether any angle formed by its adjacent edges is less than $2\pi/3$.
- 3 If there is at least one angle in G that is found to be less than $2\pi/3$, the algorithm adds a new **Fermat point** and assigns it as **balanced point**.

Big picture

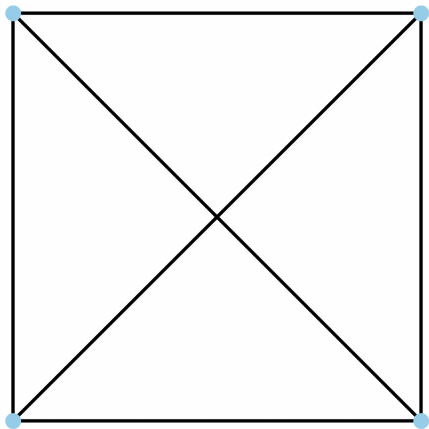
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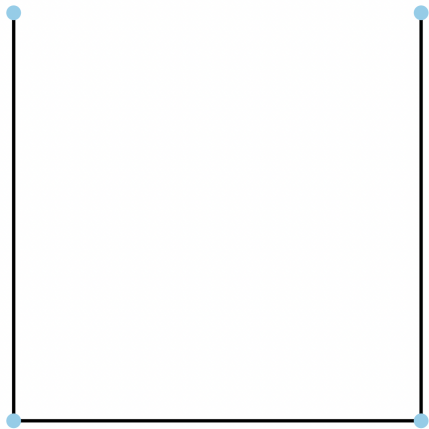
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If each unbalanced point v_u has degree 1 or connects to two other points with an angle greater than or equal to $2\pi/3$, the algorithm terminates. Otherwise, the process continues by adding a balanced point.

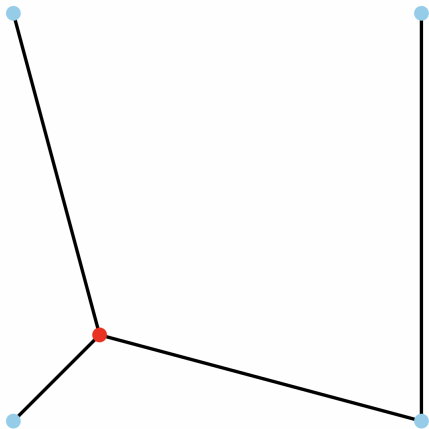
Fully Connected Graph



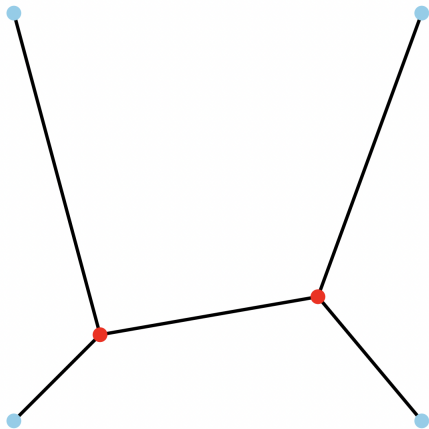
Minimum-Spanning-Tree



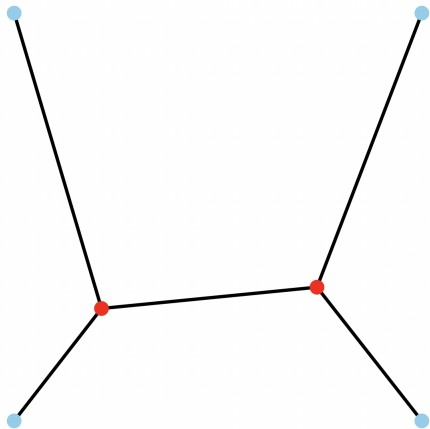
Add a new Fermat Point



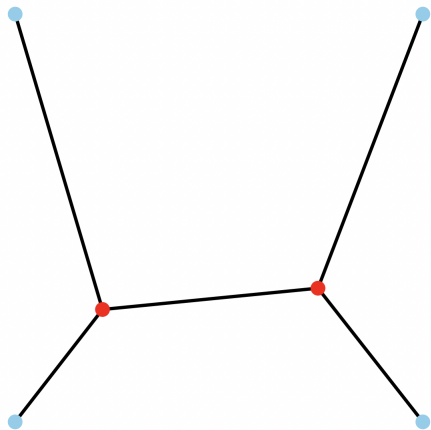
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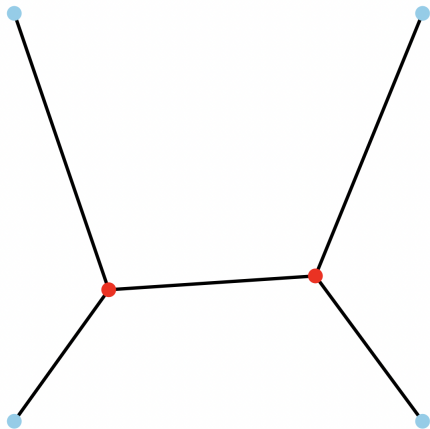
Balancing



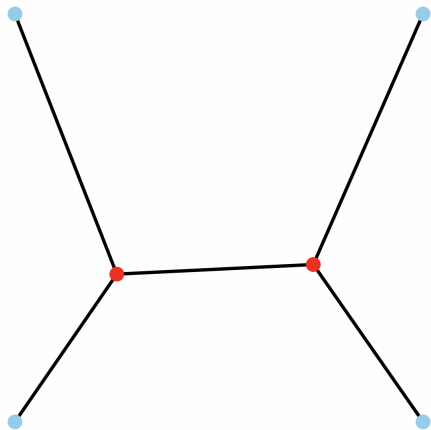
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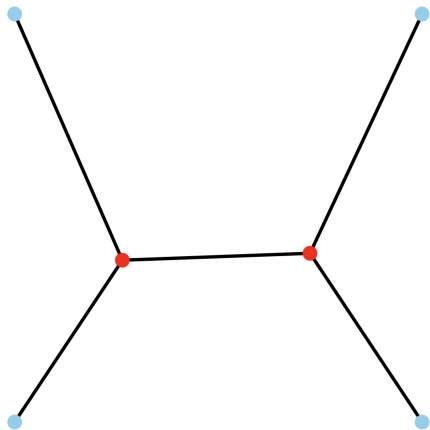
Balancing



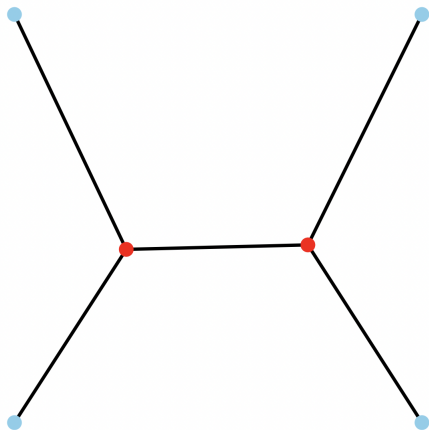
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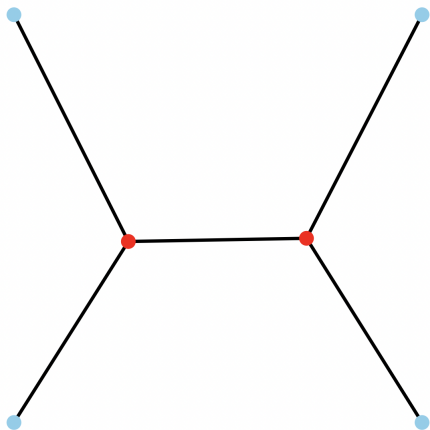
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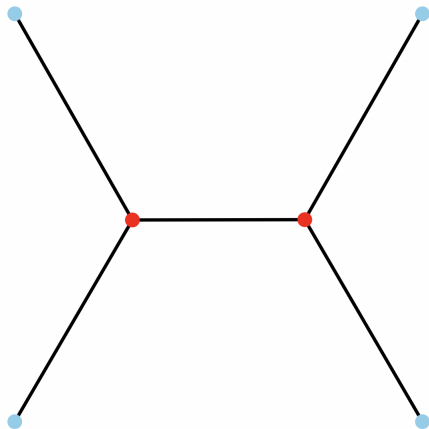
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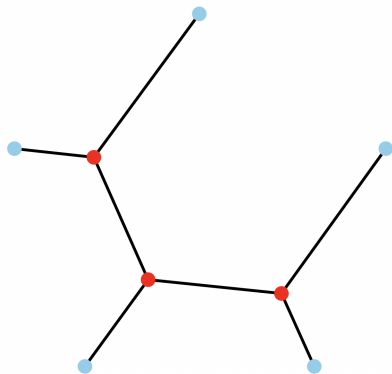
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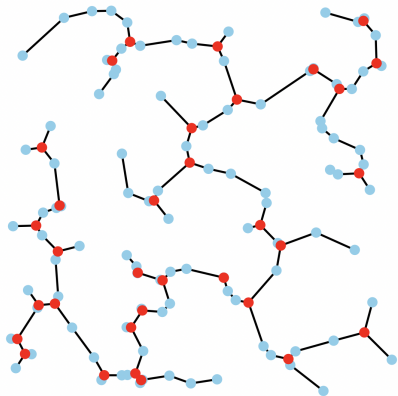
Final Result



Results



(a) 5 unbalanced points

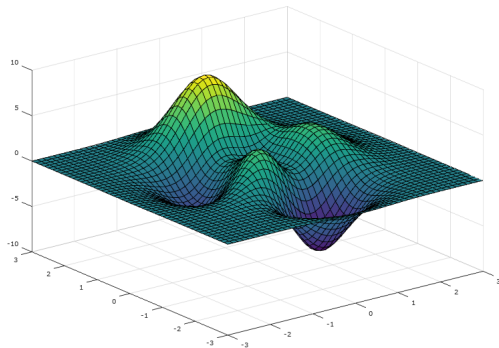


(b) 100 unbalanced points

Motivation

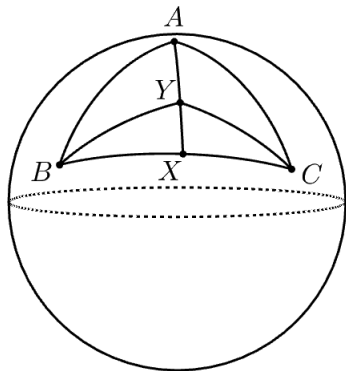
Recall that the main idea of the algorithm is adding new **balanced points**.

Question: Do **balanced points** exist on a general surface?



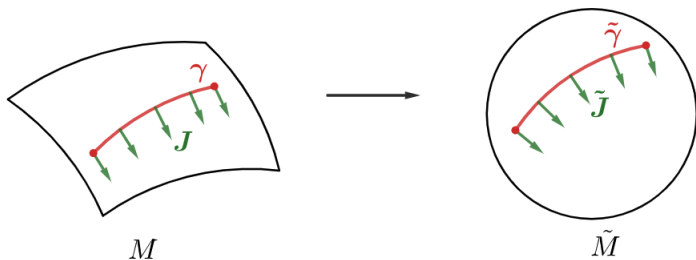
Existence of a balanced vertex - Sphere

Theorem 1 (*Sphere*). Given triangle ABC on a round sphere M with radius R such that its **three angles measure less than $2\pi/3$** . If the **maximum geodesic distance** of two points in the domain of the triangle ABC is **less than $R\pi/2$** , then **there exists a balanced point**.

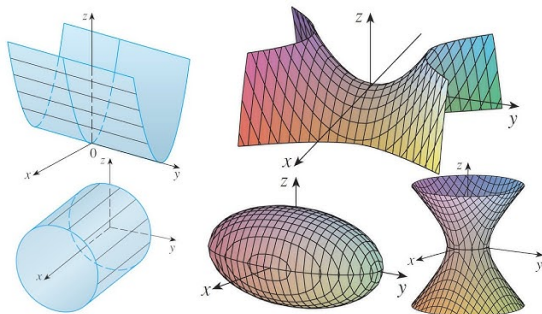


Existence of a balanced vertex - General surface

Theorem 2 (*General surface*). Let M be a Riemannian surface such that its Gaussian curvature K is **bounded above** by $1/R^2$, for $R > 0$. Let triangle ABC on M be given such that its three angles measure less than $2\pi/3$. If the maximum geodesic distance of two points in the domain of the triangle ABC is less than $R\pi/2$, then **there exists a balanced point**.



- Implement the algorithm on **different surfaces** with **different Riemannian metrics**
- Investigate the **complexity** of the algorithm





Nadja Betzler.

Steiner tree problems in the analysis of biological networks.
Masters thesis, 2006.



Michael Herring.

The euclidean steiner tree problem.
Denison Univ., Granville, OH, 2004.



Alexander Nabutovsky and Fabian Parsch.

Geodesic nets: Some examples and open problems.
Experimental Mathematics, 32(1):1–25, 2023.



(a) Thesis



(b) Paper

Nguyen, Duc Toan. "On the existence of a balanced vertex in geodesic nets with three boundary vertices." *Journal of Geometry* 116.3 (2025): 36.